

Reversible and Reversible Complement Cyclic Codes over Galois Rings

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Abstract: Let $\mathfrak{R}$ be a Galois ring of characteristic p^a where p is a prime and a is a natural number. In this paper a relation between reversible cyclic codes and reversible complement cyclic codes has been determined. The reversible complement cyclic codes have applications in DNA-based computations.

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1. Introduction

The Galois ring $\mathfrak{R} = GR(p^a, m)$ is a Galois extension of Z_{p^a} of degree m , where p is a prime and a is a natural number. The Galois ring of characteristic p^a and cardinality p^{am} , where m is a natural number, is isomorphic to the ring $Z_{p^a}[u]/f(u)$ where $f(u)$ is a monic basic irreducible polynomial of degree m in $Z_{p^a}[u]$ ^{1,2}.

A cyclic code over a ring R is a linear code which is invariant under cyclic shifts. The cyclic codes of length n over a ring R under a linear map can be considered as ideals of $R[x]/\langle x^n - 1 \rangle$. Thus the tuple $c = (c_0, c_1, \dots, c_{n-1})$ can be identified with the polynomial $c(x) = c_0 + c_1x + \dots + c_{n-1}x^{n-1}$. Moreover, the reversible complement cyclic codes have wide applications in the field of DNA-based computations. For reference see³⁻⁸.

This paper is organized as follows. In Section 2, a necessary and sufficient condition for a reversible cyclic code to be a reversible complement cyclic code has been obtained, thereby extending a result of

Bennenni³ to Galois rings. In Section 3, examples are given to illustrate the result obtained in Section 2. We have identified reversible complement cyclic codes in the list given by Abualrub⁹ for single generator reversible cyclic codes of length 4 over Z_4 .

2. Reversible and Reversible Complement Cyclic Codes

Let $\mathfrak{R} = GR(p^a, m) \cong Z_{p^a}[u]/\langle f(u) \rangle$ be a Galois ring and $\mathfrak{R}_n = \mathfrak{R}[x]/\langle x^n - 1 \rangle$.

Remark 1: It can be easily seen that for every element e in $\mathfrak{R}$ there exists an element $\bar{e}$ such that $e + \bar{e} = 1 + u + u^2 + \dots + u^{m-1}$. We shall call $\bar{e}$ the complement of e . Clearly $\overline{(\bar{e})} = e$.

Definition 1: A cyclic code C of length n over a ring R is called reversible if $(c_{n-1}, \dots, c_1, c_0)$ belongs to C whenever $(c_0, c_1, \dots, c_{n-1}) \in C$.

Definition 2: A cyclic code C of length n over a ring R is called reversible complement code if $(\bar{c}_{n-1}, \dots, \bar{c}_1, \bar{c}_0) \in C$ whenever $(c_0, c_1, \dots, c_{n-1}) \in C$.

The element $(\bar{c}_{n-1}, \dots, \bar{c}_1, \bar{c}_0)$, represented as a polynomial $(c(x))^{RC}$, is called reversible complement of $c(x)$.

Definition 3: The reciprocal polynomial of $f(x) \in R[x]$ denoted as $f^*(x)$ is defined as $f^*(x) = x^{\deg(f(x))} f(1/x)$.

Remark 2: It is easy to prove that $((c^*(x))^{RC})^* = (c(x))^{RC}$ for $c(x)$ in C .

Lemma 1: Let C be a reversible complement cyclic code over $\mathfrak{R}$ of length n . Then C is reversible.

Proof: Let $g(u) = 1 + u + u^2 + \dots + u^{m-1}$ and $0(x)$ be the zero polynomial in $\mathfrak{R}_n$. Then $0(x) \in C$ and hence $(0(x))^{RC} \in C$.

i.e. $g(u)(1 + \dots + x^{n-2} + x^{n-1}) \in C$.

Let $c(x) = c_0 + c_1x + \dots + c_tx^t, 0 \leq t < n$, be a polynomial in C . Then

$$(c(x))^{RC} = g(u)(1 + x + \dots + x^{n-t-2}) + (\bar{c}_tx^{n-t-1} + \dots + \bar{c}_1x^{n-2} + \bar{c}_0x^{n-1})$$

belongs to C . Therefore, $g(u)(1+x+\dots+x^{n-1})-(c(x))^{RC} \in C$.

Moreover,

$$\begin{aligned} g(u)(1+x+\dots+x^{n-1})-(c(x))^{RC} &= (c_t x^{n-t-1} + \dots + c_1 x^{n-2} + c_0 x^{n-1}) \\ &= x^{n-t-1}(c_t + \dots + c_1 x^{t-1} + c_0 x^t) \end{aligned}$$

Therefore, $x^{n-t-1}(c_t + \dots + c_1 x^{t-1} + c_0 x^t) \in C$.

As C is cyclic, $(c_t + \dots + c_1 x^{t-1} + c_0 x^t) \in C$.

So, $c^*(x) \in C$. Hence C is a reversible cyclic code.

Lemma 2: Let C be a reversible cyclic code over $\mathfrak{R}$ of length n such that $(1+u+u^2+\dots+u^{m-1})(1+x+x^2+\dots+x^{n-2}+x^{n-1}) \in C$. Then C is a reversible complement cyclic code.

Proof: Let $c(x) = c_0 + c_1 x + \dots + c_t x^t$, $0 \leq t < n$, be a polynomial in C . As C is cyclic,

$$x^{n-t-1}(c(x)) = c_0 x^{n-t-1} + \dots + c_{t-1} x^{n-2} + c_t x^{n-1} \in C.$$

Moreover, $g(u)(1+x+\dots+x^{n-2}+x^{n-1}) \in C$

where $g(u) = 1+u+u^2+\dots+u^{m-1}$.

Therefore, $g(u)(1+x+\dots+x^{n-2}+x^{n-1})-(c_0 x^{n-t-1} + \dots + c_{t-1} x^{n-2} + c_t x^{n-1}) \in C$.

i.e. $g(u)(1+x+\dots+x^{n-t-2})+(\bar{c}_0 x^{n-t-1} + \dots + \bar{c}_t x^{n-1}) = (c^*(x))^{RC} \in C$.

As C is reversible, $((c^*(x))^{RC})^* \in C$. This together with Remark 2 implies that $(c(x))^{RC} \in C$. Hence C is a reversible complement cyclic code.

Lemmas 1 and 2 combine to give the following theorem:

Theorem 1: A cyclic code C is a reversible complement cyclic code of length n over $\mathfrak{R}$ if and only if C is a reversible cyclic code of length n over $\mathfrak{R}$ and $(1+u+u^2+\dots+u^{m-1})(1+x+\dots+x^{n-1}) \in C$.

3. Illustration

The single generator reversible cyclic codes over $Z_4 = GR(2^2, 1)$ of length 4 as ideals of $\mathfrak{R}_4 = GR(2^2, 1)/\langle x^4 - 1 \rangle$ have been obtained by Abualrub¹. With the help of following examples, we illustrate that the

class of reversible complement cyclic codes over Z_4 of length 4 is a proper subset of the class of reversible cyclic codes over Z_4 of length 4.

Example 1: The reversible cyclic code $C = \langle x^2 + 2x + 3 \rangle$ of length 4 over $GR(2^2, 1)$ is reversible complement and $(x^3 + x^2 + x + 1) \in C$.

Example 2: The reversible cyclic code $C = \langle x^2 + 3 \rangle$ of length 4 over $GR(2^2, 1)$ is not reversible complement cyclic code and $(x^3 + x^2 + x + 1) \notin C$.

In Table 1, we provide the complete list of single generator reversible cyclic codes and reversible complement cyclic codes of length 4 over $GR(2^2, 1)$.

Table 1: Single generator reversible and reversible complement cyclic codes of length over $GR(2^2, 1)$.

Sr. No	Reversible	Reversible Complement
1.	$\langle 1 \rangle$	Yes
2.	$\langle x + 1 \rangle$	Yes
3.	$\langle x + 3 \rangle$	Yes
4.	$\langle x^2 + 1 \rangle$	Yes
5.	$\langle x^2 + 3 \rangle$	No
6.	$\langle x^2 + 2x + 1 \rangle$	No
7.	$\langle x^2 + 2x + 3 \rangle$	Yes
8.	$\langle x^3 + x^2 + x + 1 \rangle$	Yes
9.	$\langle x^3 + 3x^2 + x + 3 \rangle$	No
10.	$\langle 2x^3 + 2x^2 + 2x + 2 \rangle$	No
11.	$\langle 2x^2 + 2 \rangle$	No
12.	$\langle 2x + 2 \rangle$	No
13.	$\langle 2 \rangle$	No

4. Conclusion

In this paper a necessary and sufficient condition for a reversible cyclic code to be a reversible complement cyclic code has been obtained.

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