A New Goal Programming Approach for Multi-Objective Solid Transportation Problem with Interval-Valued Intuitionistic Fuzzy Logic

M. K. Sharma, Kamini and Nitesh Dhiman

Department of Mathematics C.C.S. University, Meerut-250004, India Email: drmukeshsharma@gmail.com, Kaminivarun123@gmail.com, niteshdhiman91@gmail.com

Vandana

Dept of Computer Applications, SCRIET, C.C.S. University, Meerut Email: ranavd@gmail.com

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Abstract: Many researchers considered fuzzy and intuitionistic fuzzy parameters in the transportation problem, but we deal with intervalvalued intuitionistic fuzzy parameters, which is another type of uncertainty that covers the favourable as well as unfavourable cases. In a multi-objective transportation problem, it is difficult to find the optimal solution for all objectives simultaneously. So, we have used the Intuitionistic Fuzzy Goal Programming (IFGP) with deviational function $d_{\nu} = (1 - w_{\nu})$ (where w_{ν} is fixed numerical weight for d_{ν} decided the importance of highest acceptance level which of k^{th} objective relative to other objectives) to find the compromise optimal solution. In this paper, we extend the model proposed by Nomani M.A. et al. in 2016 for Multi-Objective Solid Transportation Problem (MOSTP) with interval-valued intuitionistic fuzzy cost. We applied this approach to find the solution of the Solid Transportation Problem (STP) with satisfying all the constraints. To find the compromise optimal solution for all objectives simultaneously, we apply a new IFGP approach. The main focus of the proposed approach is to minimize the all objectives simultaneously and to obtain the solution nearly closed to the lower bounds of objectives. A numerical example is being carried out in favour of the proposed algorithm. Keywords: Multi-Objective Transportation Problem (MOTP), Solid Transportation Problem (STP), Interval-Valued Intuitionistic Fuzzy Numbers (IVIFN), Goal Programming (GP), Intuitionistic Fuzzy Goal Programming(IFGP). **2010 AMS Classification Number:** 90C08, 90C70

1. Introduction

Transportation problem is one of the linear programming problem, but it is a special type problem due to its structure. So it can not be solved by the usual simplex method. The main objective of transportation problem is to minimize the transport cost by satisfiving the demand of all costumers. Hitchcock¹ in 1941 introduced the concept of Classical Transportation Problem. Transportation cost depends on several conditions. In real life situation the parameters of the problem consists of various type of uncertainties. To deal with such uncertainties, Zadeh² in 1965 introduced the concept of fuzzy logic that helps to represent such type of imprecise data and Bellman³ in 1970 used this theory in decision making problems. In current situation, if the transportation problem contains three kinds of constraints such as demand, supply and conveyance then these types of transportation problems are known as a Solid Transportation Problem (STP) introduced by Haley⁴ in 1962. Classical transportation problem is a special type of Solid Transportation Problem when the number of conveyance in solid transportation problem is only one.

Transport cost is not only the objective of transportation problem but it also compromises to deal with more than one objective. There are more than one objectives of transportation problem which are conflicting with each other. Such type of transportation problems possessing this property are known as Multi-Objective Transportation Problems (MOTP). To find an optimal solution for all objectives of such type of problem is tedious task. Zimmermann⁵ introduced fuzzy programming and linear programming with multi objectives and also⁶ used fuzzy set theory in mathematical programming. Bit⁷ et.al. gave a method to solve the multi-objective transportation problem with interval parameters. Then in 1961, Charnes and Cooper⁸ introduced a powerful approach to deal with multi-objective known as "Goal Programming", which is very helpful to find acompromise optimal solution for MOTP. It is also known as the generalization of linear programming problem dealing with multi-objectives. Narasimhan⁹ introduced the Goal Programming approach in uncertain or fuzzy environment. Radhakrishnan and Anukokila¹⁰ introduced the Fractional Goal Programming for solving Fuzzy Solid Transportation Problem with interval cost. Pal B.B. & Moitra¹¹ used Fuzzy Goal Programming (FGP) for Long Range Production Planning in Agricultural Systems by using fixed numerical weight $w_k = \frac{1}{U_k - L_k}$. El-Wahed and Lee¹² solved the MOTP by using Interactive Fuzzy Goal Programming. Also, Zangiabadi and Maleki¹³ used FGP for solving MOTP. There are many approaches for solving multiobjective problems with fuzzy and intuitionistic fuzzy parameters, but Atanassov¹⁴ and Atanasson and Gargov¹⁵ introduced interval-valued intuitionistic fuzzy set in which membership and non-membership degree of any element is in interval form. Aggarwal and Gupta¹⁶ used a new ranking method based on signed distance for solving Intuitionistic Fuzzy Solid Transportation Problem. In this research article we extended the model proposed by Nomani M.A. et. al.¹⁷ for Multi-Objective Solid Transportation Problem (MOSTP) with interval-valued intuitionistic fuzzy cost.

The research paper is divided into eight sections. In the first section of the research paper deals with the introductory part of the problem and in the second section, we defined some basic concepts which are being used to our problem. The third section of the paper is related to deffuzzification of the Interval Valued Intuitionistic Fuzzy Numbers known as the the accuracy function for IVIFN. Fourth section of the paper we formulated the problem mathematically. We proposed our model in section fifth and the algorithm for getting the compromise solution of the multi-objective problems in the sixth section of the research paper. In the seventh section of the research paper a numerical is being carried out in support of our proposed model and algorithm to show the efficiency of our proposed model. In the last section of our research paper we gave the conclusion and interpretation of the results.

2. Preliminaries

2.1. Fuzzy Set: Let *X* be any universal set and \overline{A} be any subset of *X* then \overline{A} is fuzzy set if it is denoted by $\overline{A} = \{x, \mu_{\overline{A}}(x)\}$, where $x \in X, \mu_{\overline{A}}(x) \colon X \to [0,1]$ and $\mu_{\overline{A}}(x)$ is membership grade of x in \overline{A} .

2.2. Fuzzy Number: A fuzzy set \overline{A} on the real line is called a fuzzy number if it satisfies the following conditions;

(a) \overline{A} is normal i.e. there exist any $x \in \overline{A}$ such that $\mu_{\overline{A}}(x) = 1$.

(b) \overline{A} is convex set i.e.

 $\mu_{\overline{A}}(\lambda x_1 + (1 - \lambda) x_2) \ge \min(\mu_{\overline{A}}(x_1), \mu_{\overline{A}}(x_2)), \ \forall x_1, x_2 \in \overline{A}, \lambda \in [0, 1].$

(c) Membership function is piecewise continuous.

2.3. Intuitionistic Fuzzy Set: Let X be any universal set and \overline{A} be any subset of X then \overline{A} is intuitionistic fuzzy set if it is denoted by $\overline{A} = \{x, \mu_{\overline{A}}(x), \nu_{\overline{A}}(x)\}$, where $x \in X$, $\mu_{\overline{A}}(x) \colon X \to [0,1]$, $\nu_{\overline{A}}(x) \colon X \to [0,1]$

 $\mu_{\bar{A}}(x) \& v_{\bar{A}}(x)$ denotes membership grade & non-membership grade of x in \bar{A} with hesitation part, $\pi_{\bar{A}}(x) = 1 - (\mu_{\bar{A}}(x) + v_{\bar{A}}(x))$, where $0 \le \pi_{\bar{A}}(x) \le 1$.

2.4. Intuitionistic Fuzzy Number: An Intuitionistic fuzzy set \overline{A} on the real line is called an intuitionistic fuzzy number if it satisfies following conditions;

(a) \overline{A} is normal i.e. there exist any $x \in \overline{A}$ such that $\mu_{\overline{A}}(x) = 1$ or $\nu_{\overline{A}}(x) = 0$.

(b) \overline{A} is convex set for membership function

$$\mu_{\overline{A}}(\lambda x_1 + (1 - \lambda) x_2) \ge \min(\mu_{\overline{A}}(x_1), \mu_{\overline{A}}(x_2)), \forall x_1, x_2 \in \overline{A}, \ \lambda \in [0, 1]^{-1}$$

(c) \overline{A} is concave set for non-membership function

 $v_{\overline{A}}(\lambda x_1 + (1 - \lambda) x_2) \le \max(v_{\overline{A}}(x_1), v_{\overline{A}}(x_2)), \forall x_1, x_2 \in \overline{A}, \lambda \in [0, 1]$

(d) Membership and non-membership functions are piece-wise continuous.

2.5. Interval-Valued Intuitionistic Fuzzy Number (IVIFN): An interval-valued intuitionistic fuzzy number \overline{A} is denoted by

$$\overline{A} = \{(a_1^L, a_3^U, a_2, a_4^L, a_5^U), (b_1^L, b_3^U, a_2, b_4^L, b_5^U)\}$$

with the lower and upper membership and non-membership is as follows



Figure 1.

Mathematically, the membership function and non-memebrship function for lower as well as upper bound will be formulated as follows:

2.5.1 Lower membership function: Lower bound for membership function is defined as

$$\mu_{\overline{A}}^{L}(x) = \begin{cases} 1, & x = a_{2}, \\ \frac{x - a_{1}^{L}}{a_{2} - a_{1}^{L}}, & a_{1}^{L} < x < a_{2}, \\ \frac{a_{4}^{L} - x}{a_{4}^{L} - a_{2}}, & a_{2} < x < a_{4}^{L}, \\ 0, & a_{1}^{L} \ge x, x \ge a_{4}^{L}. \end{cases}$$

2.5.2. Upper membership function: Upper bound for membership function is defined as

$$\mu_{\overline{A}}^{U}(x) = \begin{cases} 1, & x = a_{2}, \\ \frac{x - a_{3}^{U}}{a_{2} - a_{3}^{U}}, & a_{3}^{U} < x < a_{2}, \\ \frac{a_{5}^{U} - x}{a_{5}^{U} - a_{2}}, & a_{2} < x < a_{5}^{U}, \\ 0, & a_{3}^{U} \ge x, x \ge a_{5}^{U}. \end{cases}$$

2.5.3 Lower non-membership function: Lower bound non-membership function is defined as

$$v_{\bar{A}}^{L} = \begin{cases} 0, & x = a_{2}, \\ \frac{a_{2} - x}{a_{2} - b_{1}^{L}}, & b_{1}^{L} < x < a_{2}, \\ \frac{x - a_{2}}{b_{4}^{L} - a_{2}}, & a_{2} < x < b_{4}^{L}, \\ 1, & x \le b_{1}^{L}, x \ge b_{4}^{L}. \end{cases}$$

2.5.4. Upper non-membership function: Upper bound for non-membership function is defined as

$$v_{\bar{A}}^{U}(x) = \begin{cases} 0, & x = a_{2}, \\ \frac{a_{2} - x}{a_{2} - b_{3}^{U}}, & b_{3}^{U} < x < a_{2}, \\ \frac{x - a_{2}}{b_{5}^{U} - a_{2}}, & a_{2} < x < b_{5}^{U}, \\ 1, & x \le b_{3}^{U}, x \ge b_{5}^{U}, \end{cases}$$

3. Defuzzification (Accuracy) Function for Interval Valued Intuitionistic Fuzzy Number

Let $\overline{A} = \{x, (a_1^L, a_3^U, a_2, a_4^L, a_5^U), (b_1^L, b_3^U, a_2, b_4^L, b_5^U)\}$ be an IVIFN. Then these inte- rval valued intuotionistic fuzzy numbers will be converted into closed inte- rvals using $\alpha - cut$ method as follows:

$$u^{L}(\alpha) = [a_{1}^{L} + \alpha(a_{2} - a_{1}^{L}), a_{4}^{L} - \alpha(a_{4}^{L} - a_{2})],$$

$$u^{U}(\alpha) = [a_{3}^{U} + \alpha(a_{2} - a_{3}^{U}), a_{5}^{U} - \alpha(a_{5}^{U} - a_{2})],$$

$$v^{L}(\alpha) = [b_{1}^{L} - (1 - \alpha)(a_{2} - b_{1}^{L}), a_{2} + (1 - \alpha)(b_{4}^{L} - a_{2})],$$

$$v^{U}(\alpha) = [b_{3}^{U} - (1 - \alpha)(a_{2} - b_{3}^{U}), a_{2} + (1 - \alpha)(b_{5}^{U} - a_{2})],$$

where $0 < \alpha < 1$.

Then by using parametric form of IVIFN, Helipern¹⁸ introduced the defuzzified(accuracy) value for IVIFN. Which is as follows:

$$R(\overline{A}) = \frac{a_1^L + a_3^U + b_1^L + b_3^U + 8 * a_2 + a_4^L + a_5^U + b_4^L + b_5^U}{16}$$

4. Mathematical Formulation of the Problem

4.1. Multi-Objective Transportation Problem (MOTP): Let there be *m* sources and *n* destination in the transportation problem and a_i be available quantity of items at i^{th} sources and b_j be the demand of j^{th}

destination. Let x_{ij} be the number of items which have to be transported from i^{th} source to j^{th} destination. Let c_{ij} be the cost per unit of the item which has to be transported from i^{th} source to j^{th} destination, where i=1,2....m and j=1,2....n. The aim of the problem is to find the quantity of items x_{ij} at which the all objectives of problem are optimized.

Let there be K number of objectives in transportation problem, then the mathematical formulation of the problem is:

$$MinZ_k(x) = \sum_{i=1}^m \sum_{j=1}^n c_{ij}^k x_{ij}$$
, where $k = 1, 2....K$.

Subject to the constraints:

$$\sum_{j=1}^{n} x_{ij} = a_i \text{, where } i = 1, 2, \dots, m, \sum_{i=1}^{m} x_{ij} = b_j \text{, where } j = 1, 2, \dots, n$$
$$\sum_{i=1}^{m} a_i = \sum_{j=1}^{n} b_j \text{ where } i = 1, 2, \dots, m, j = 1, 2, \dots, n, x_{ij} \ge 0.$$

4.2. Multi-Objective Solid Transportation Problem (MOSTP): Let there be *m* sources, *n* destination and *D* no. of conveyance in the problem. Let a_{jd} be available quantity of items which transported to j^{th} destination by d^{th} conveyance, b_{id} be the demand of the product which transported from i^{th} source by d^{th} conveyance and e_{ij} be the quantity of product which transported from i^{th} source to j^{th} destination. Let x_{ijd} be the no. of products transport from i^{th} source to j^{th} destination by d^{th} conveyance and c_{ijd} be the cost of per unit of product which transport from i^{th} source to j^{th} destination by d^{th} conveyance and c_{ijd} be the cost of per unit of product which transport from i^{th} source to j^{th} destination by d^{th} conveyance to j^{th} destination by d^{th} conveyance and c_{ijd} be the cost of per unit of product which transport from i^{th} source to j^{th} destination by d^{th} conveyance to j^{th} destination by d^{th} conveyance and c_{ijd} be the cost of per unit of product which transport from i^{th} source to j^{th} destination d^{th} conveyance, where i=1,2...,m, j=1,2...,n and d=1,2...,D. Then the mathematical formulation of MOSTP:

$$MinZ_k(x) = \sum_{i=1}^{m} \sum_{j=1}^{n} \sum_{d=1}^{D} c_{ijd}^k x_{ijd}$$
, where $k = 1, 2, ..., K$.

Subject to the constraints:

$$\sum_{i=1}^{m} x_{ijd} = a_{jd} , \sum_{j=1}^{n} x_{ijd} = b_{id} , \sum_{d=1}^{D} x_{ijd} = e_{ij}$$

where,

$$\sum_{j=1}^{n} a_{jd} = \sum_{i=1}^{m} b_{id} , \qquad \sum_{j=1}^{n} e_{ij} = \sum_{d=1}^{D} b_{id} , \qquad \sum_{i=1}^{m} b_{ij} = \sum_{d=1}^{D} a_{jd} ,$$
$$\sum_{j=1}^{n} \sum_{d=1}^{D} a_{jd} = \sum_{d=1}^{D} \sum_{i=1}^{m} b_{di} = \sum_{i=1}^{m} \sum_{j=1}^{n} a_{ij} , \ x_{ijd} \ge 0 ,$$

where i = 1, 2, ..., m, j = 1, 2, ..., n and d = 1, 2, ..., D.

4.3. Goal Programming Problem (GPP): Goal Programming was introduced by Charnes and Cooper in 1961. The aim of GP is to minimize the deviational of objectives Z_k from aspiration levels \overline{Z}_k of the objectives which are to be decided by decision maker. When the given problem is of minimizing nature then the over achievement deviational variable d_k^+ will be minimized whenever problem is of maximizing nature then under achievement deviational variable d_k^- will be maximized, where the mathematical form of the over achievement deviational variable is of the form

$$d_{k}^{+} = \frac{1}{2} \{ (Z_{k} - \overline{Z}_{k}) + |Z_{k} - \overline{Z}_{k}| \}, Z_{k} \ge \overline{Z}_{k}$$

and the mathematical form of the under achievement deviational variable is of the form

$$d_{k}^{-} = \frac{1}{2} \{ (\overline{Z}_{k} - Z_{k}) + |\overline{Z}_{k} - Z_{k}| \}, \overline{Z}_{k} \ge Z_{k}$$

The mathematical formulation of goal programming:

$$\operatorname{Min} Z_k(x) = \sum_{i=1}^m \sum_{j=1}^n \sum_{d=1}^p c_{ijd}^k x_{ijd} - \overline{Z}_k \quad .$$

Subject to $\sum_{i=1}^{m} x_{i}$

$$x_{ijd} = a_{jd}$$
, $\sum_{j=1}^{n} x_{ijd} = b_{id}$, $\sum_{d=1}^{D} x_{ijd} = e_{ij}$,

where,

$$\sum_{j=1}^{n} a_{jd} = \sum_{i=1}^{m} b_{id} , \sum_{j=1}^{n} e_{ij} = \sum_{d=1}^{D} b_{id} , \sum_{i=1}^{m} b_{ij} = \sum_{d=1}^{D} a_{jd} ,$$

$$\sum_{j=1}^{n} \sum_{d=1}^{D} a_{jd} = \sum_{d=1}^{D} \sum_{i=1}^{m} b_{di} = \sum_{i=1}^{m} \sum_{j=1}^{n} a_{ij} , \ x_{ijd} \ge 0 ,$$

where i = 1, 2, ..., m, j = 1, 2, ..., n and d = 1, 2, ..., D.

Where Z_k and \overline{Z}_k are k^{th} objective function and aspiration level of k^{th} objective function respectively. Now, the mathematical formulation is as follows:

$$\operatorname{Min}\sum_{k=1}^{K}(d_{k}^{+}-d_{k}^{-})$$

Subject to

$$\sum_{i=1}^{m} \sum_{j=1}^{n} \sum_{d=1}^{p} c_{ijd}^{k} x_{ijd} - \overline{Z}_{k} = d_{k}^{+} - d_{k}^{-}$$

$$\sum_{i=1}^{m} x_{ijd} = a_{jd} , \quad \sum_{j=1}^{n} x_{ijd} = b_{id} , \quad \sum_{d=1}^{D} x_{ijd} = e_{ij} , d_{k}^{+} - d_{k}^{-} \ge 0, \ k = 1, 2...K$$

4.3.1. Fuzzy Goal Programming (FGP): Let L_k and U_k be lower and upper bounds of Z_k objective function and let $\mu_k(Z_k)$ be a membership function for Z_k objective function in minimizing problem. We consider $\mu_k(Z_k)$ function as linear function with the highest membership grade 1, which defines as;

$$\mu_{k}(Z_{k}(x)) = \begin{cases} 1, & \text{if } Z_{k} \leq L_{k}, \\ \frac{U_{k} - Z_{k}(x)}{U_{k} - L_{k}} & \text{if } L_{k} \leq Z_{k}(x) \leq U_{k}, \\ 0, & \text{if } Z_{k} \geq U_{k}. \end{cases}$$

By applying min-max approach, mathematical model of fuzzy goal programming;

 $Min\theta$,

Subject to

$$\mu_k(Z_k(x)) - d_k^+ + d_k^- = 1$$

$$\sum_{i=1}^m x_{ijd} = a_{jd} , \sum_{j=1}^n x_{ijd} = b_{id} , \sum_{d=1}^D x_{ijd} = e_{ij}$$

$$\theta \ge d_k^-, k = 1, 2...K, d_k^+ d_k^- = 0, \ \theta \ge 0, \theta \le 1, x_{iid} \ge 0,$$

where i = 1, 2, ..., m, j = 1, 2, ..., n and d = 1, 2, ..., D.

Now, if DM is not satisfied then choose new upper or lower bounds of the objectives and again apply same process.

5. Proposed Approach for the MOTP Model

For MOTP, there are several approaches to solve the problem and find the solution for all objectives simultaneously. In our proposed method, we used the new fuzzy goal programming approach with the deviational function $d_k = (1 - w_k)$ in the place of deviational variables d_k^- and d_k^+ where $w_k = \frac{1}{U_k - L_k}$ is fixed numerical weight associated with d_k (under achevement deviational variable) which decided the importance of the highest acceptance level of k^{th} objective relative to other objectives. In the existing paper¹³, the model is used to find the compromise optimal solution for Multi-Objective Transportation Problem (MOTP) with crisp data, we extend it for Multi-Objective Solid Transportation Problem (MOSTP) with Interval-Valued Intuitionistic Fuzzy Cost (IVIFC) and also take underachievement deviational variable d_{k} for membership function of each objective, used fixed $w_k = \frac{1}{U_k - L_k}$ associated to each d_k . By using the linear membership function with deviational function $d_k = (1 - w_k)$, the membership goal with highest acceptance of membership grade 1 is represented as:

$$\mu_k(z_k(x)) + d_k(1 - w_k) = 1,$$

We know that highest acceptance of membership goal is 1, there exist only underachievement deviational variable d_k . Because, if overachievement deviational variable exist that means the aspiration level of fuzzy goal is more than 1 which is not possible. So only under achievement deviational variable is required for achieving the highest aspired level for membership goal. Now, by applying the min-max approach of goal programming, the proposed model is as follows;

Min
$$\theta$$

 $\mu_k(Z_k(x)) + d_k(1 - w_k) = 1, \ \theta \ge d_k(1 - w_k), \ k = 1, 2....K$

$$\sum_{i=1}^{m} x_{ijd} = a_{jd} , \sum_{j=1}^{n} x_{ijd} = b_{id} , \sum_{d=1}^{D} x_{ijd} = e_{ij} ,$$
$$x_{ijd} \ge 0 , d_{k} \ge 0 , 0 \le \theta \le 1$$

where,

$$\mu_{k}(z_{k}(x)) = \begin{cases} 1, & \text{if } Z_{k} \leq L_{k}, \\ \frac{U_{k} - Z_{k}(x)}{U_{k} - L_{k}}, & \text{if } L_{k} \leq Z_{k}(x) \leq U_{k}, \\ 0, & \text{if } Z_{k} \geq U_{k}. \end{cases}$$

linear function for minimize problem.

6. Proposed Algorithm for MOSTP

Step-1: First, Consider a Multi-Objective Solid Transportation Problem (MOSTP) with Interval-Valued Intuitionistic Fuzzy Cost (IVIFC).

Step-2: Convert the problem into crisp form by using defuzzification (accuracy) function.

Step-3: Check whether the given problem is balanced otherwise convert into balanced problem by adding a dummy row or Column with zero cost.

Step-4: Now, solve the problem by Vogel's Approximation Method taken single objective at a time for each j respectively. Then, we get basic feasible solutions for all objectives of the problem.

Step-5: Now, apply the method¹⁶ proposed to interval valued intuitionistic fuzzy cost to find the optimal solution for each objective. The method will be as follows:

- (i) After find the solution by vogel's approximation method, for each j=1, check allocation of c_{i1d} exceeds b_{id} or not. If not exceeds then find the solution for each j by vogel's approximation method.
- (ii) If exceeds by θ units, then making a loop starting with c_{i1d} and move horizontally and vertically with allocation filled up cells. Now, assign $+\theta$ and $-\theta$ sign alternatively at each corner started $-\theta$ at c_{i1d} . All cells with must be filled up with minimum θ allocation.
- (iii) Now, if any entry at c_{i1d} with +ve sign in loop is not filled up, then minimum value at b_{id} must be θ . If it is filled up with β value then

minimum value at b_{id} must be $\theta + \beta$. Do the same process till all three constraints are not satisfied.

Step-6: Now, make the payoff matrix for all objectives as shown in the table 1, where X_1, X_2, \dots, X_k be the Solution obtained by step-5 for Z_1, Z_2, \dots, Z_k respectively then;

	X_{1}	X_{2}		X_k
Z_1				
Z_2				
•	•	•	•	•
•	•	•		•
Z_k				

Table 1. Payoff matrix for all objectives

Step-7: Determine the lower bound L_k and upper bound U_k for k^{th} objective function given in the problem.

Step-8: Now, to get a compromise optimal solution for all objectives simultaneously, built the model for given problem then solve by using Lingo software¹⁹.

7. Numerical Example

We consider a Multi-objective solid transportation problem as shown in table with two objectives.

	D_1	D_2	D_3	Supply
<i>O</i> ₁	$e_{11} = 0$ $\overline{c}_{111}, \overline{c}_{112}, \overline{c}_{113}$	$e_{12} = 6$ $\overline{c}_{121}, \overline{c}_{122}, \overline{c}_{123}$	$e_{13} = 9$ $\overline{c}_{131}, \overline{c}_{132}, \overline{c}_{133}$	<i>b</i> ₁₁ <i>b</i> ₁₂ <i>b</i> ₁₃ 6, 9, 10
<i>O</i> ₂	$e_{21} = 21$ $\overline{c}_{211}, \overline{c}_{212}, \overline{c}_{213}$	$e_{22} = 9$ $\overline{c}_{221}, \overline{c}_{222}, \overline{c}_{223}$	$e_{23} = 14$ $\overline{c}_{231}, \overline{c}_{232}, \overline{c}_{233}$	$b_{21}b_{22}b_{23} \\ 13, 14, 17$

 Table 2. Solid Transportation Problem

<i>O</i> ₃	$e_{31} = 21$ $\overline{c}_{311}, \overline{c}_{312}, \overline{c}_{313}$	$e_{32} = 13$ $\overline{c}_{321}, \overline{c}_{322}, \overline{c}_{323}$	$e_{33} = 12$ $\overline{c}_{331}, \overline{c}_{332}, \overline{c}_{333}$	$\begin{array}{c} b_{31} b_{32} b_{33} \\ 15, 13, 18 \end{array}$
Demand	<i>a</i> ₁₁ <i>a</i> ₁₂ <i>a</i> ₁₃ 15, 17, 20	<i>a</i> ₂₁ <i>a</i> ₂₂ <i>a</i> ₂₃ 8, 11, 9	$a_{31} a_{32} a_{33}$ 11, 8, 16	

Now we will apply our proposed algorithm for this numerical problem **Step 1& 2:**

Where for First objective:

$$\begin{split} \overline{c}_{111} &= (5,7,9,11,13)(4,5,9,13,14), & R(\overline{c}_{111}) = 9; \\ \overline{c}_{112} &= (4,6,7,8,9)(3,4,7,9,10), & R(\overline{c}_{112}) = 6.8; \\ \overline{c}_{113} &= (6,7,9,11,12)(5,6,9,12,13), & R(\overline{c}_{113}) = 9; \\ \overline{c}_{211} &= (8,9,10,12,13)(7,8,10,13,14), & R(\overline{c}_{211}) = 10.3; \\ \overline{c}_{212} &= (3,5,6,7,9)(2,3,6,9,10), & R(\overline{c}_{212}) = 6; \\ \overline{c}_{213} &= (2,3,5,7,8)(1,2,5,8,9), & R(\overline{c}_{213}) = 5; \\ \overline{c}_{311} &= (5,6,7,8,9)(3,4,7,10,11), & R(\overline{c}_{311}) = 7; \\ \overline{c}_{312} &= (2,3,5,7,8)(1,2,5,8,9), & R(\overline{c}_{312}) = 5; \\ \overline{c}_{313} &= (3,4,5,6,7)(1,2,5,7,8), & R(\overline{c}_{312}) = 5; \\ \overline{c}_{122} &= (8,9,10,12,13)(7,8,10,13,14), & R(\overline{c}_{122}) = 10.3; \\ \overline{c}_{123} &= (2,3,5,7,8)(1,2,5,8,9), & R(\overline{c}_{123}) = 5; \\ \overline{c}_{221} &= (3,5,6,7,9)(2,3,6,9,10), & R(\overline{c}_{221}) = 6; \\ \overline{c}_{222} &= (5,7,9,11,13)(4,5,9,13,14), & R(\overline{c}_{222}) = 9; \\ \overline{c}_{223} &= (7,8,10,11,12)(6,7,10,12,13), & R(\overline{c}_{321}) = 5.8; \\ \overline{c}_{322} &= (10,11,12,13,14)(8,9,12,15,16), & R(\overline{c}_{322}) = 12; \\ \overline{c}_{323} &= (6,7,10,11,12)(4,5,10,12,13), & R(\overline{c}_{323}) = 9; \\ \overline{c}_{131} &= (10,11,12,13,14)(8,9,12,15,16), & R(\overline{c}_{131}) = 12; \\ \overline{c}_{132} &= (8,9,10,12,13)(7,8,10,13,14), & R(\overline{c}_{132}) = 10.3; \\ \overline{c}_{133} &= (6,7,10,11,12)(4,5,10,12,13), & R(\overline{c}_{132}) = 10.3; \\ \overline{c}_{133} &= (6,7,10,11,12)(4,5,10,12,13), & R(\overline{c}_{132}) = 10.3; \\ \overline{c}_{133} &= (6,7,10,11,12)(4,5,10,12,13), & R(\overline{c}_{133}) = 9; \\ \overline{c}_{133} &= (6,7,10,11,12)(4,5,10,12,13), & R(\overline{c}_{1$$

$$\begin{split} \overline{c}_{231} &= (3,4,5,6,7)(1,2,5,7,8), & R(\overline{c}_{231}) = 4.8; \\ \overline{c}_{232} &= (5,6,7,8,9)(3,4,7,10,11), & R(\overline{c}_{232}) = 7; \\ \overline{c}_{233} &= (2,3,5,7,8)(1,2,5,8,9), & R(\overline{c}_{233}) = 5; \\ \overline{c}_{331} &= (6,7,10,11,12)(4,5,10,12,13), & R(\overline{c}_{331}) = 9; \\ \overline{c}_{332} &= (5,7,9,11,13)(4,5,9,13,14), & R(\overline{c}_{332}) = 9; \\ \overline{c}_{333} &= (5,6,7,8,9)(3,4,7,10,11), & R(\overline{c}_{333}) = 7; \end{split}$$

And for second objective:

$$\begin{split} \overline{c}_{111} &= (5,7,9,11,13)(4,5,9,13,14), & R(\overline{c}_{111}) = 9; \\ \overline{c}_{112} &= (10,11,12,13,14)(8,9,12,15,16), & R(\overline{c}_{112}) = 12; \\ \overline{c}_{113} &= (8,9,10,12,13)(6,7,10,14,15), & R(\overline{c}_{113}) = 10.3; \\ \overline{c}_{211} &= (8,9,10,12,13)(6,7,10,14,15), & R(\overline{c}_{211}) = 10.3; \\ \overline{c}_{212} &= (6,7,9,11,12)(5,6,9,12,13), & R(\overline{c}_{212}) = 9; \\ \overline{c}_{213} &= (4,6,7,8,9)(3,4,7,9,10), & R(\overline{c}_{213}) = 6.8; \\ \overline{c}_{311} &= (4,6,7,8,9)(3,4,7,9,10), & R(\overline{c}_{311}) = 6.8; \\ \overline{c}_{312} &= (6,7,9,11,12)(5,6,9,12,13), & R(\overline{c}_{312}) = 9; \\ \overline{c}_{313} &= (2,3,5,7,8)(1,2,5,8,9), & R(\overline{c}_{313}) = 5; \\ \overline{c}_{122} &= (6,7,9,11,12)(5,6,9,12,13), & R(\overline{c}_{122}) = 9; \\ \overline{c}_{123} &= (4,5,7,8,9)(2,3,7,10,11), & R(\overline{c}_{122}) = 9; \\ \overline{c}_{223} &= (7,8,10,11,12)(6,7,10,12,12), & R(\overline{c}_{223}) = 6.8; \\ \overline{c}_{222} &= (10,11,12,13,14)(8,9,12,15,16), & R(\overline{c}_{222}) = 12; \\ \overline{c}_{223} &= (8,9,10,12,13)(6,7,10,14,15), & R(\overline{c}_{322}) = 10.3; \\ \overline{c}_{322} &= (7,8,10,11,12)(6,7,10,12,13), & R(\overline{c}_{323}) = 9.8; \\ \overline{c}_{323} &= (7,8,10,11,12)(6,7,10,12,13), & R(\overline{c}_{323}) = 9.8; \\ \overline{c}_{333} &= (5,7,9,11,13)(4,5,9,13,14), & R(\overline{c}_{133}) = 9; \\ \overline{c}_{333} &= (5,7,9,11,13)(4,5,9,13,14), & R(\overline{c}_{233}) = 9.8; \\ \overline{c}_{233} &= (7,8,10,11,12)(5,6,9,12,13), & R(\overline{c}_{333}) = 9; \\ \overline{c}_{333} &= (5,7,9,11,13)(4,5,9,13,14), & R(\overline{c}_{233}) = 9.8; \\ \overline{c}_{233} &= (7,8,10,11,12)(5,7,10,12,13), & R(\overline{c}_{233}) = 9.8; \\ \overline{c}_{233} &= (7,8,10,11,12)(5,6,9,12,13), & R(\overline{c}_{233}) = 9.8; \\ \overline{c}_{233} &= (7,8,10,11,12)(5,7,10,12,13), & R(\overline{c}_{233}) = 9.8; \\ \overline{c}_{333} &= (3,5,6,7,9)(2,3,6,9,10), & R(\overline{c}_{331}) = 6; \\ \end{array}$$

$$\overline{c}_{332} = (2,3,5,7,8)(1,2,5,8,9), \qquad R(\overline{c}_{332}) = 5;$$

$$\overline{c}_{333} = (3,5,6,7,9)(2,3,6,9,10), \qquad R(\overline{c}_{333}) = 6;$$

Step 3: Since

$$\sum_{j=1}^{n} a_{jd} = \sum_{i=1}^{m} b_{id} , \qquad \sum_{j=1}^{n} e_{ij} = \sum_{d=1}^{D} b_{id} , \qquad \sum_{i=1}^{m} b_{ij} = \sum_{d=1}^{D} a_{jd} ,$$
$$\sum_{j=1}^{n} \sum_{d=1}^{D} a_{jd} = \sum_{d=1}^{D} \sum_{i=1}^{m} b_{di} = \sum_{i=1}^{m} \sum_{j=1}^{n} a_{ij} ,$$

Thus, problem is balanced.

Step 4 & 5: Now, by applying step 4 and 5 we get the optimal solution for both objectives such as:

For first objective solution X_1 as follows:

$$x_{112} = 7, x_{113} = 3, x_{212} = 4, x_{213} = 17, x_{311} = 15, x_{312} = 6, x_{121} = 1, x_{122} = 2, x_{123} = 3,$$

$$x_{221} = 7, x_{222} = 2, x_{322} = 7, x_{323} = 6, x_{131} = 5, x_{133} = 4, x_{231} = 6, x_{232} = 8, x_{333} = 12.$$

For second objective solution X_2 as follows:

$$x_{111} = 6, x_{112} = 4, x_{211} = 4, x_{212} = 13, x_{213} = 4, x_{311} = 5, x_{313} = 16, x_{122} = 5,$$

$$x_{123} = 1, x_{222} = 1, x_{223} = 8, x_{321} = 8, x_{322} = 5, x_{133} = 9, x_{231} = 9, x_{233} = 5,$$

$$x_{331} = 2, x_{332} = 8, x_{333} = 2.$$

Step 6: Make a payoff matrix as shown in table 3.

 Table 3. Payoff matrix

	X_1	X_2
Z_1	822	835.7
Z_2	950	899.7

Step 7: Now, $L_1 = 822, U_1 = 835.7; L_2 = 899.7, U_2 = 950$, where L_1 and U_1 are lower and upper bounds of Z_1 (first objective function) and L_2 , U_2 are lower and upper bounds of Z_2 respectively (second objective function). **Step 8:** Built Model for the MOSTP:

Min θ ,
$\mu_1(Z_1(x)) + d_1(1 - w_1) = 1$; where $w_1 = 0.073$,
$Z_1(x) - 12.7d_1 = 822;$
$\mu_2(Z_2(x)) + d_2(1 - w_2) = 1$; where $w_2 = 0.02$,
$Z_2(x) - 49.29d_2 = 899.7;$
$\theta \ge 0.927 d_1$, $\theta \ge 0.980 d_2$

$$\begin{split} x_{111} + x_{112} + x_{113} &= 10; \ x_{211} + x_{212} + x_{213} &= 21; \ x_{311} + x_{312} + x_{313} &= 21; \\ x_{111} + x_{211} + x_{311} &= 15; \ x_{112} + x_{212} + x_{312} &= 17; \ x_{113} + x_{213} + x_{313} &= 20; \\ x_{121} + x_{122} + x_{123} &= 6; \ x_{221} + x_{222} + x_{223} &= 9; \ x_{321} + x_{322} + x_{323} &= 13; \\ x_{121} + x_{221} + x_{321} &= 8; \ x_{122} + x_{222} + x_{322} &= 11; \ x_{123} + x_{223} + x_{323} &= 9; \\ x_{131} + x_{132} + x_{133} &= 9; \ x_{231} + x_{232} + x_{233} &= 14; \ x_{331} + x_{332} + x_{333} &= 12; \\ x_{131} + x_{231} + x_{331} &= 11; \ x_{132} + x_{232} + x_{332} &= 8; \ x_{133} + x_{233} + x_{333} &= 16; \\ x_{111} + x_{121} + x_{131} &= 6; \ x_{112} + x_{122} + x_{132} &= 9; \ x_{113} + x_{123} + x_{133} &= 10; \\ x_{211} + x_{221} + x_{231} &= 13; \ x_{212} + x_{222} + x_{232} &= 14; \ x_{213} + x_{223} + x_{233} &= 17; \\ x_{311} + x_{321} + x_{331} &= 15; \ x_{312} + x_{322} + x_{332} &= 13; \ x_{313} + x_{323} + x_{333} &= 18; \\ 0 \leq d_1 \leq 1, 0 \leq d_2 \leq 1, 0 \leq \theta \leq 1, \end{split}$$

Then, global solution of the problem obtained by Lingo 18:

 $\begin{aligned} x_{111} = 6, & x_{112} = 4, x_{211} = 1.696429, x_{212} = 8.151786, x_{213} = 11.151786, x_{311} = 7.303571, \\ x_{312} = 4.848214, x_{313} = 8.848214, x_{112} = 2.848214, x_{123} = 3.151786, x_{221} = 3.151786, \\ x_{223} = 5.848214, x_{321} = 4.848214, x_{322} = 8.151786, x_{132} = 2.151786, x_{133} = 6.848214, \\ x_{231} = 8.151786, x_{232} = 5.848214, x_{331} = 2.848214, x_{333} = 9.151786. \\ Z_1 = 822 \text{ and } Z_2 = 909.5. \end{aligned}$

i.e.,

i.e.,

8. Conclusion

We applied new approach to find the compromise solution of the Multi-Objective Transportation Problem (MOTP). We take special transportation problem known as the Solid Transportation Problem (STP) in which three constraints are included. We use Lingo software to find the compromise optimal solution of the problem. It is a new approach to find the global solution for Multi-Objective Solid Transportation Problem (MOSTP) with Interval-Valued Intuitionistic Fuzzy Cost (IVIFC). In this approach, we take a different deviational function $d_{k} = (1 - w_{k})$ where w_{k} is fixed and $w_{k} = 1/U_{k} - L_{k}$ associated to each d_{k} . By this proposed method, we found best solution for multi-objective solid transportation problem. We compared our solution with Fuzzy Goal Programming (FGP) approach. Our solution is more optimal than the solution obtained by FGP with linear membership function. The solution obtained by FGP $Z_1 = 835.7$ and $Z_2 = 947.39$ which is more than $Z_1 = 822$ and $Z_2 = 909.5$ solution obtained by our proposed algorithm. This proves that our proposed algorithm is better than the previous algorithm.

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